# Computer Networks

Gnkgo, Informatik B. Sc. 4. Semester

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### 1 Internet

- Internet connects End Systems/ Hosts by a system of communications links and package switches
- Packets: data segment and header
- IP (Internet Protocol): Specifies the format for packets sent among routers and end systems
  - unreliable service
  - best-effort delivery
  - modify before forwarding the packet to the next hop
    - \* destination MAC address
    - \* Time-to-live
    - \* IP checksum
- DSL (digital subscriber line): broadband residential access
  - downstream data channel: tens to few hundred Mbps
  - upstream data channel: few Mbps to few tens Mbps
- FDM (Frequency-division multiplexing): link dedicates a frequency band on the band to a communication for the duration of connection. The width of this band is called bandwidth.
- TDM (time-division multiplexing): time is divided into frames of fixes size, and each frame is split into a fixed amount of slots. These slots are dedicated to one connection

## 2 Packet, Circuit Switching

- Packet switching relies on buffers to account for unexpected bursts
- Circuit switching is better applicable when the peak-to-average utilization ratio is low.
- large Peak / Average  $\rightarrow$  On-demand
- small peak /average  $\rightarrow$  circuit

### 3 Delays

- nodal processing delay:
  - Time used to decode the header, determine which output needs to be used, bit-level checking, etc.
  - Negligible in comparison to other delays, little room for improvement here
- queuing delay:
  - Time that the packet has to wait in the queue until the desired output will be free
  - Better load balancing (i.e., via routing), reduce queue size
- transmission delay:
  - Time after the packet gets out of the queue until it's on the output line; let L be the length of the packet, R the transmission rate from router A to B, then the transmission delay is L/R.
  - Increase the link capacity
- propagation delay:
  - Time the package takes to propagate through the physical medium (e.g., at light speed with radio waves),
     i.e., it takes on the way from router A to B
  - If possible, use data replication to shorten the client-server distance. Improving routing or physical transmission media also helps

## 4 Layers



- Application Layer: Application protocols and layers reside here. The internet includes many protocols here, as HTTP (provides Web document requests and responses), SMTP (for e-mail transfer), and FTP (file transfer from host to host). The DNS is also an application layer protocol. Packet of information at this layer is called a message.
- **Transport Layer:** Transports application layer messages between application endpoints. In the internet, there's TCP and UDP. E.g. longer messages are split into shorter segments. Packets in this layer are called segments (TCP & UDP) or Datagram (UDP).
- Network Layer: The layer is responsible for moving datagrams from host to host. It receives a segment and an address. This layer includes the IP, which defines the fields in a datagram and how to work on those fields. Packets in this layer are called datagrams / packet.
  - ARP: broadcast domain receives everyone until a router.
  - Goes through MAC.
- Link Layer: The network layer brings a datagram from one node to another, but to move a packet from one specific node to another, it relies on the link layer.
  - It delivers the packet to the next node. E.g. Ethernet or WiFi. Link layer protocols are called frames.
  - Error detection
- An Ethernet switch can interconnect a 10 Mbps Ethernet network and a 1 Gbps Ethernet network.
- The 802.11b wireless protocol incorporates a link-layer ACK not present in regular Ethernet.
- **Physical layer:** The job of the link layer is to move whole frames from one node to another; however, the job of the physical layer is to move individual bits. Many protocols exist, depending on the physical medium; e.g. Ethernet has different protocols for different cable types.

## 5 Address Resolution Protocol (ARP)

• It allows a host to get the MAC address associated with an IP address.

### 6 Ethernet

- 1. Frame starts with destination address.
  - If not destined to this device, it can be simply dropped.
- 2. Source address.
- 3. Type / length field.
  - Tells receiver what payload this frame carries.
- 4. Data and pad.
  - Minimum length of an Ethernet frame is 64 bytes in total.
- 5. Frame check sequence.
  - Cyclic Redundancy Check (CRC) for error checking.

## 7 Wireless Local Area Networks (WLANs)

- Multiple access points can operate on different channels.
- Nodes are often mobile, requiring dynamic routing.

# 8 TCP/IP Model

| TCP/IP Model            |  |
|-------------------------|--|
| Application Layer       |  |
| Transport Layer         |  |
| Internet Layer          |  |
| Network Interface Layer |  |

• The network interface layer encompasses both the link and physical layers.

## 9 Congestion Control

- A state in which one link in a subnet is so overwhelmed that it cannot effectively handle the amount of traffic that could use it.
- A link can become congested even if it has sufficient bandwidth.
- Network layer congestion can occur even if no link is congested.

#### 9.1 Receive Procedure

1. Divide and check for zero remainder

#### 9.2 Path Lookup



Figure 1: Path Lookup

#### 9.3 Recipe to get IP addressing:

### 10 Quiz

• The Maximum Segment Size (MSS) of TCP is equal to: MSS = MTU - header(IP) - header(TCP)

| Column Name       | Explanation   |
|-------------------|---|
| Prefix            | The prefix notation (CIDR) representing the subnet's network                  |
|                   | size. IP address followed by a slash and the number of significant            |
|                   | bits in the subnet mask.  |
| # of hosts        | Maximum number of usable hosts in the subnet. Calculated as                   |
|                   | $2^{(32-\text{prefixlength})} - 2$ , accounting for network and broadcast ad- |
|                   | dresses.  |
| Prefix mask       | Subnet mask in dotted-decimal notation, each segment represents               |
|                   | 8 bits.   |
| Network Address   | Obtained by bitwise AND operation between IP address and sub-                 |
|                   | net mask.   |
| Broadcast Address | Highest address in the subnet, obtained by flipping all host bits             |
|                   | to 1s in the network address.   |
| Last Host Address | Highest usable IP address, obtained by setting all host bits to 1s            |
|                   | in the network address, except for the last bit.                              |

 Table 1: IP Addressing Recipe



Figure 2: Mac vs IP-Address

- UDP sockets type is SOCK\_DGRAM while TCP sockets type is SOCK\_STREAM.
- For a SOCK\_STREAM, an operating system stores both local and remote port.
- Given a directed graph G(V, E), with |V| and |E| being the numbers of vertices and edges, how many variables do you need for the max-flow LP formulation discussed in class?
- There is no connection establishment in UDP.
- The objective of flow control is not to overwhelm the hosts.
- The objective of congestion control is to not overwhelm the network.
- During congestion avoidance in TCP, the successful acknowledgment of a segment results in the sender congestion window growing by one segment per RTT.
- TCP (SOCK\_STREAM) is a connection-based protocol. The connection is established, and the two parties have a conversation until the connection is terminated by one of the parties or by a network error.
- UDP (SOCK\_DGRAM) is a datagram-based protocol. You send one datagram and get one reply, and then the connection terminates.

- Given a directed graph G(V, E), with |V| and |E| being the numbers of vertices and edges, how many variables do you need for the max-flow LP formulation discussed in class?: O(|E|)
- QUIC can handle switching from WiFi to a cellular network without having to reestablish the connection:
  - Uses connection IDs independent of the IP address instead of a 4-tuple like TCP to identify connections.
  - This way packets using the connection ID are still valid, even if the source IP address changes.
- If we create our simple query (dig @a.root-servers.net www.ethz.ch) and send it to the first root-server, we don't get a www.ethz.ch IP in return. Why is this?
  - Root server does not support recursive resolution.
- The minimum information one needs to resolve any DNS hostname is the IP address of a DNS root server.

## 11 Quic

- Operate in Application and Transport layer.
- Combines connection and TLS handshake  $\rightarrow$  reducing the connection setup time by one RTT.
- Enables Zero-RTT communication if the hosts have communicated before (improved handshake).
- Connection hand-over is possible by identifying connection with a connection ID instead of the 5(/4)-tuple (even with changing IP addresses e.g. when changing networks with a mobile device).
- Resolves head-of-line blocking by the logical abstraction of streams (contrary to TCP, which required you to open multiple parallel TCP connections).
- Middleboxes and NAT routers are known to drop unfamiliar transport layer protocols. Quic uses UDP to give inseparability with existing hardware.